

Definition.

Verbally. Let $p > 1$ be an integer.

p is called **prime**

if it has only two positive divisors: 1 and p .

Symbolically. Let $p \in \mathbb{Z}$ and $p > 1$.

p is **prime** $\iff \forall k \in \mathbb{Z}_+ (k \mid p \implies k = 1 \vee k = p)$.

Let $n > 1$ be an integer.

n is called **composite**

if it has more than two positive divisors.

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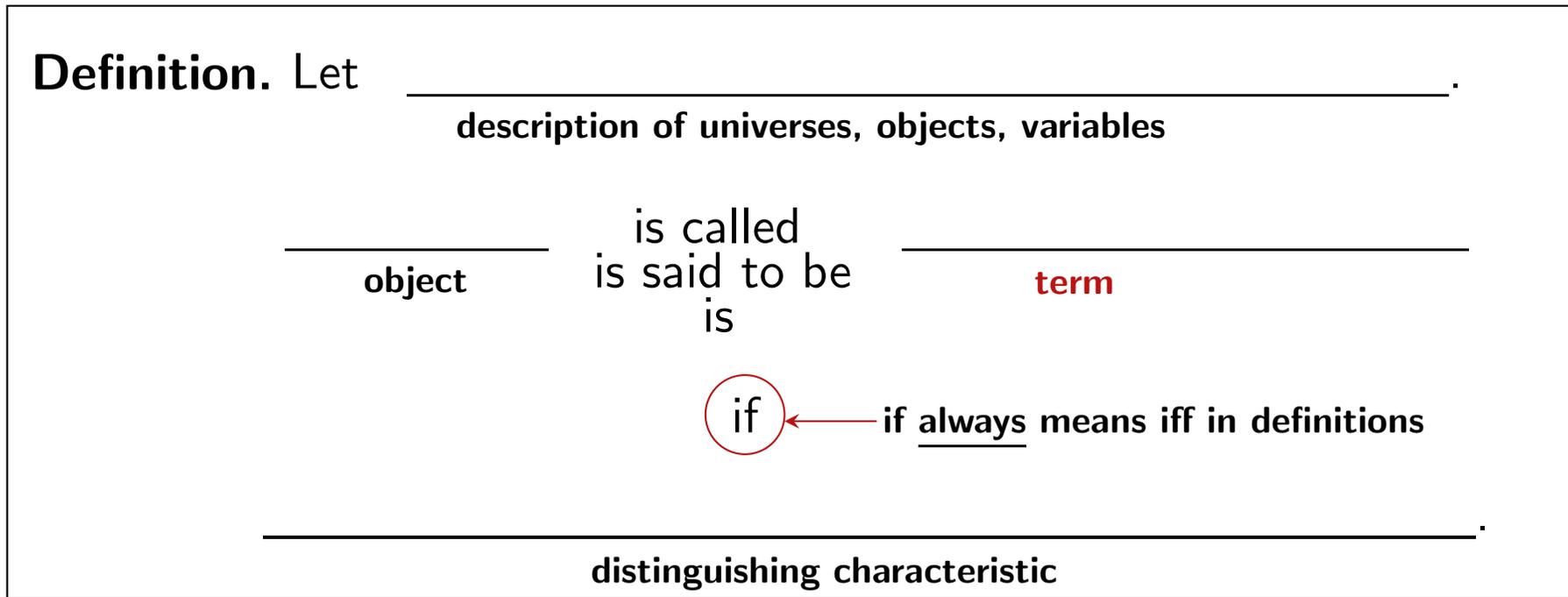
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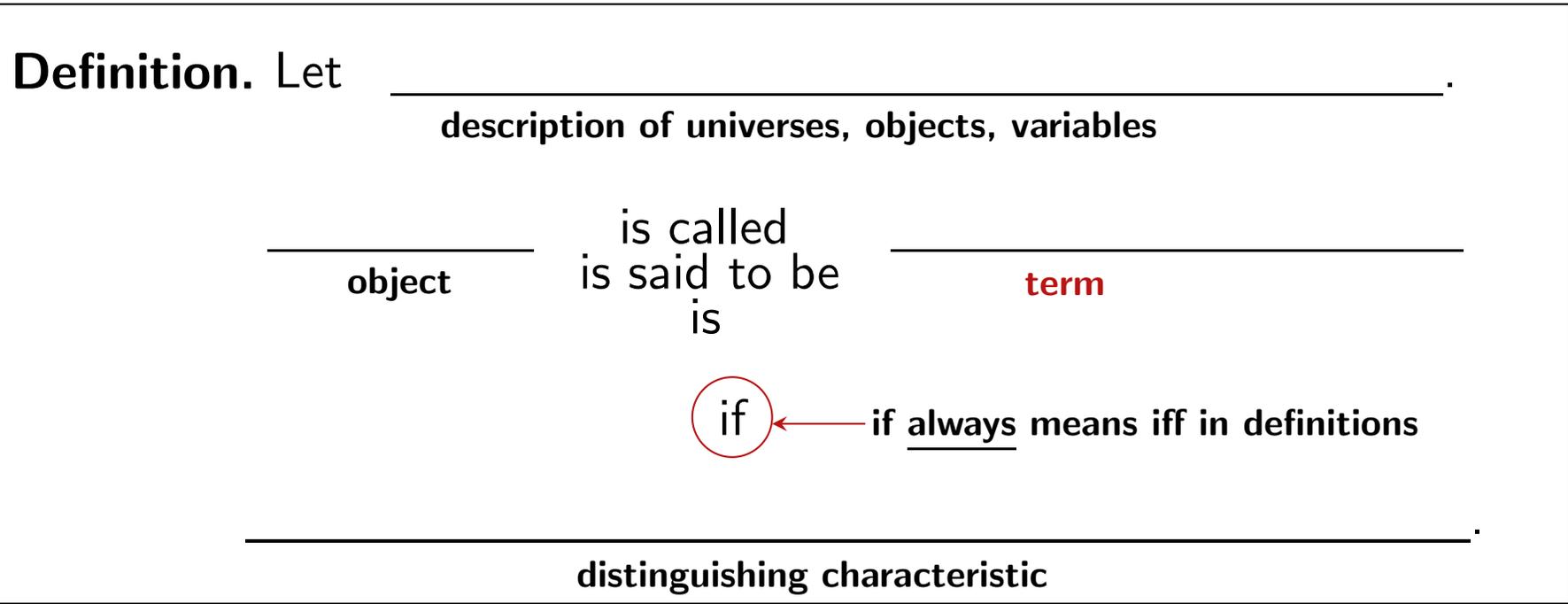
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Then $\sqrt{2} = \frac{p}{q}$ for some $p, q \in \mathbb{Z}$, $q \neq 0$.

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It means that p^2 is even.

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But p is also even, that is $2 \mid p$. We have got that $2 \mid p$ and $2 \mid q$.

$\sqrt{2}$ is irrational

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Euclid's theorem

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For source and comments see

Euclid's Elements, Book IX, Proposition 20.

How Euclid wrote this

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If possible, let it be so. Now A , B , and C measure DE , therefore G also measures DE .

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Listen to the proof and try to write it down...

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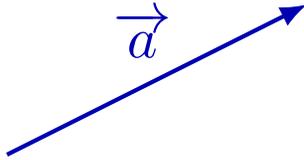
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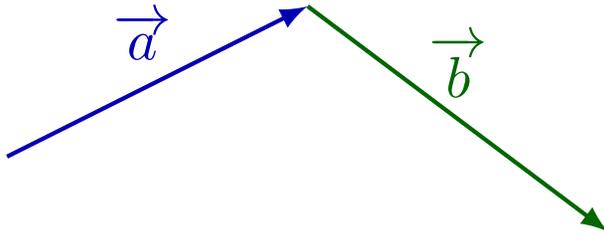
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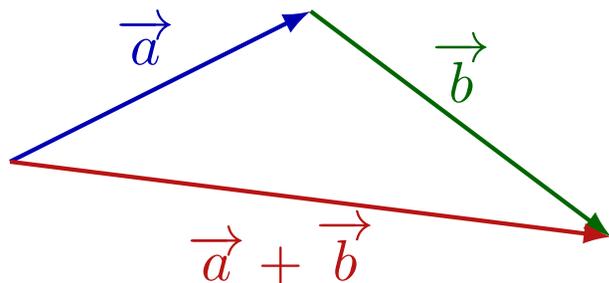
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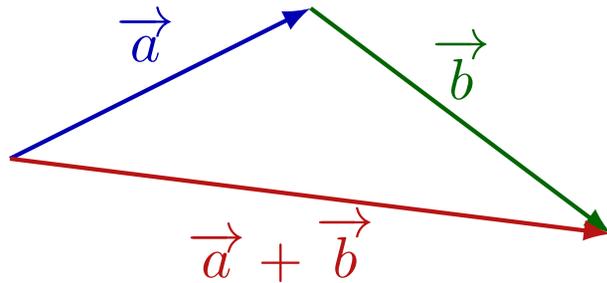
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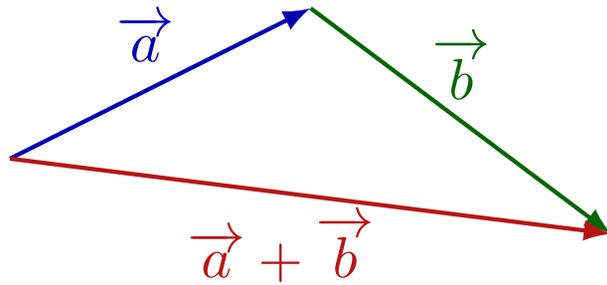
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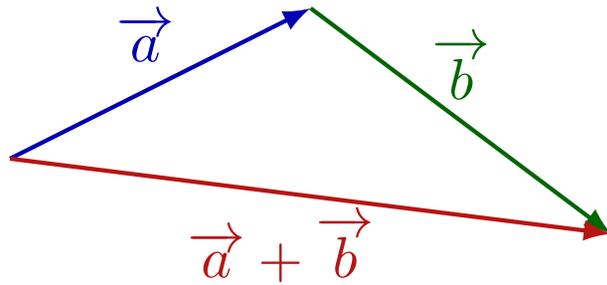


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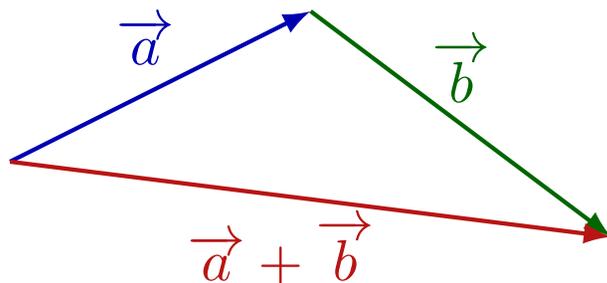
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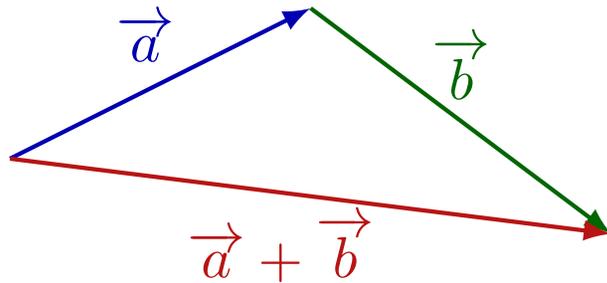
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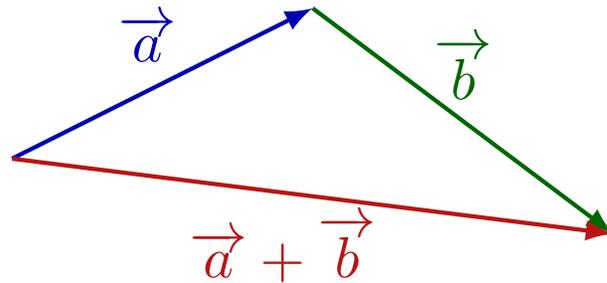
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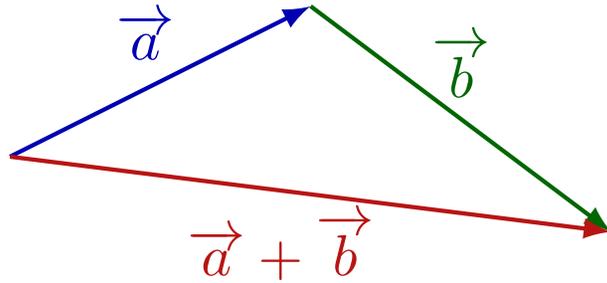
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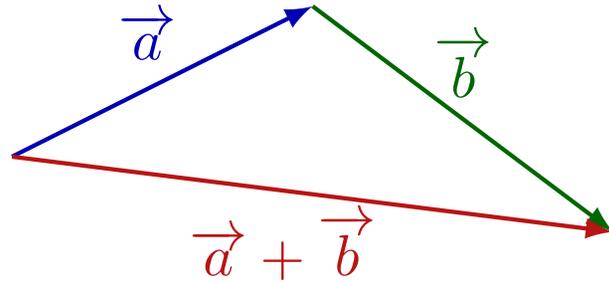
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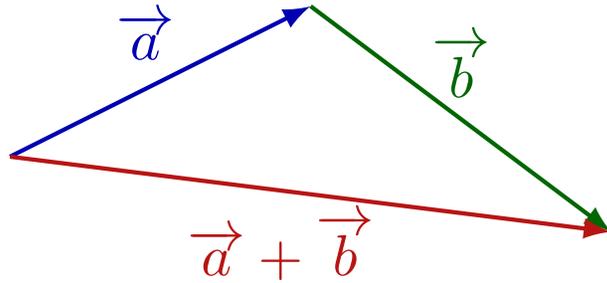
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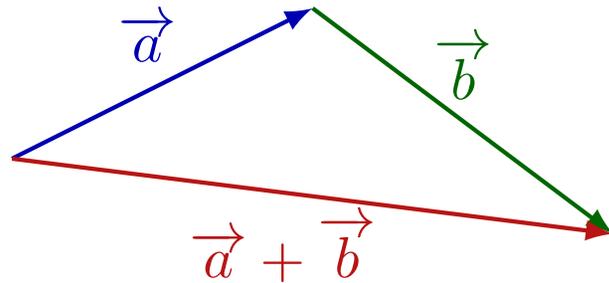
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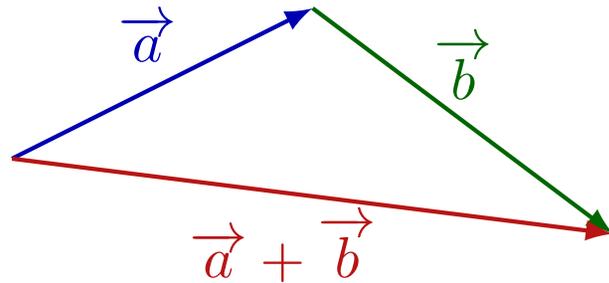
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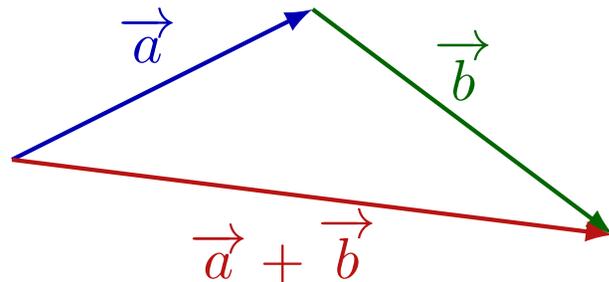
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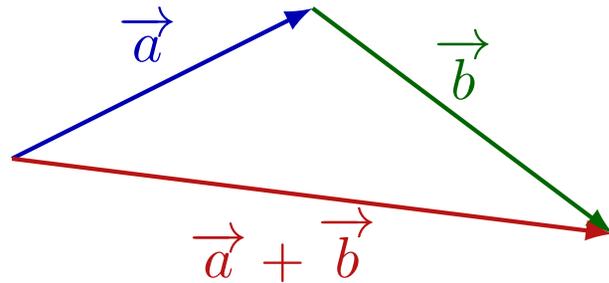
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$$|a + b| \leq ||a| + |b||.$$

Since $||a| + |b|| = |a| + |b|$,

The triangle inequality, another proof

Let us give another proof of the triangle inequality.

For any real numbers a and b , we have

$$(a + b)^2 = a^2 + b^2 + 2ab \underbrace{\leq}_{ab \leq |ab|} a^2 + b^2 + 2|ab| = |a|^2 + |b|^2 + 2|a||b| = (|a| + |b|)^2.$$

Therefore, $(a + b)^2 \leq (|a| + |b|)^2$. From this we get

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Since $||a| + |b|| = |a| + |b|$, we get $|a + b| \leq |a| + |b|$.

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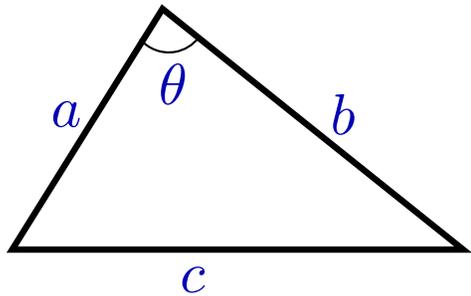
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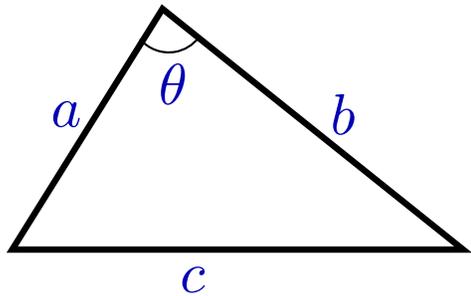
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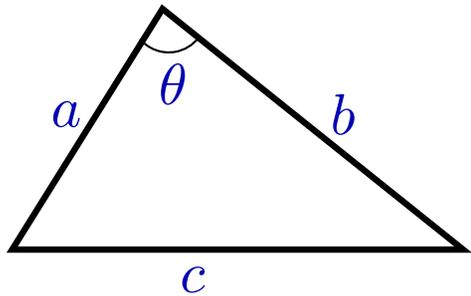
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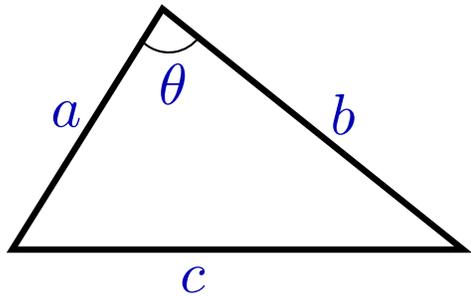
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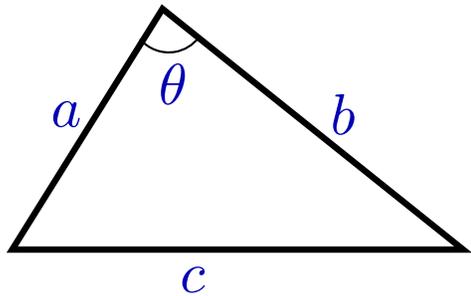
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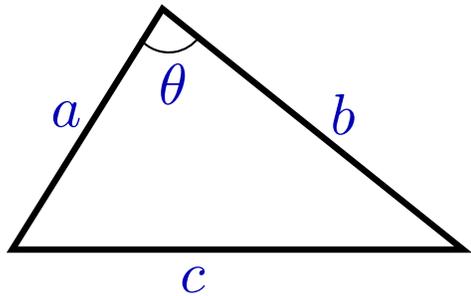
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